

Source Routing on the Edge

Scale, Reliability and Programmability for EXARINGs Internet Peering

Agenda

- 1. Who am I
- 2. State of Packet Forwarding
- 3. Requirements of modern Packet Forwarding
- 4. Solution
 - a. Data Plane
 - b. Control Plane
 - c. Issues
- 5. Questions

Who am I?

- → Oliver Herms aka takt
- → Senior Network Engineer @ EXARING AG
- → Friend of robustness, reliability, velocity
- → Network Automation Enthusiast
- → Golang and gRPC fanboy



State of Packet Forwarding

State of Packet Forwarding



Limitations of current state (1)

- → Packets to an ISP follow a single shortest path or a number of equal paths
 - All active links get the same amount of traffic
 - 10G + 100G = 20G usable capacity
 - What is equal can be tuned administratively

Limitations of current state (2)

Traffic Engineering can make use of non-shortest paths

- → Manual tweaking of Route attributes
 - Dangerous: Mistakes can cause outages
- → Only on a per Prefix basis (IP Ranges, 256-2M addresses)
- → Requires changes in Router configs
 - We fully generate them. But we review them manually.

Limitations of current state (3)

All IP Routes must be installed into Routers

- → Memory is limited
- → Expensive licenses required for 100k+ Routes
- \rightarrow Limits future growth with current platform
- → Stops us from using even cheaper Routers

Requirements

Requirements



Requirements

- → Make non-equal speed links usable
- → Make non-equal cost links usable
- → Automatically maximize utilization of cheapest links
- → Automatically move excess traffic to next cheapest link
- → Allow to take link quality into account in routing decision
- → React to changes quickly and repair any situation automatically, if possible

Nice to haves

- → Do not change Router configs
- → Support arbitrary amount of Routes
- → Allow per IP traffic engineering



Solution (1)

- → Let Vendor Routers forward traffic but not route it
 - Too inflexible to meet our needs
- → Source Routing: Let the source of traffic decide which path a packet takes
- → Servers send labeled packets
- → Packets get encapsulated into tunnels to Egress Routers

Solution (2)

- → Labeled packet arrives at Router
 - Static forwarding
 - Label indicates next-hop
 - Ignoring IP Routing Table

Advantages

- → Allows fine granular control of link utilization
 - will save € in OPEX
- → No need for IP Routing on Routers anymore
 - will save € in CAPEX)

Data Plane

Architecture Overview (Data Plane)



Dest.	Label	Tunnel	Weight
ISP A	1000	core01.fra01	10G
ISP B	2000	core01.lej01	10G
ISP B	2001	core02.lej01	20G

MPLS Label Switching Paths

- → Multiprotocol Label Switching (MPLS)
- → Label Switching Path (LSP) allows choosing Next-Hops per Label

oherms@core02.fra01> ...nces CLOSEDNET protocols mpls static-label-switched-path coffee_62_69_146_95
transit 1001899 {
 description rdev=AS201701,rif=ECIX-FRA,ndev=ECIX-FRA,nif=ECIX-FRA-001,nrole=IXP;
 next-hop 62.69.146.95;
 pop;
}

Getting to the Peering Router (PR)

- → Full MPLS deployment on internal network
 - IS-IS SR (Segment Routing)
 - LDP (Label Distribution Protocol)
 - RSVP (Resource Reservation Protocol)

- → MPLS in a Tunnel
 - MPLS over GRE/IP
 - MPLS over UDP/IP

Packet Stack leaving Machines

Tunnel IP	ader MPLS Label	IP Header of	TCP/UDP	Data
Header Port 663	5 (Next Hop)	Payload	Header	

Machines (Linux)

1. Create Foo Over UDP (FOU) encapsulated SIT tunnel per Router

modprobe fou
ip fou add port 6635 ipproto 4
ip link add name cn-cr01fra01-0 type sit remote 192.168.1.1 local
192.168.1.2 ttl 64 encap fou encap-sport 6635 encap-dport 6635

2. Add MPLS encapsulated tunnel interface routes

modprobe mpls_iptunnel
modprobe mpls_gso
ip route add 192.0.2.0/24 encap mpls 123 dev cn-cr01fra01-0

Decap MPLS-in-UDP Firewall Filter

```
oherms@core02.fra01> show configuration firewall family inet filter
CN MATROSCHKA
term MPLS-IN-UDP {
    from {
        destination-prefix-list {
            CN MATROSCHKA CORE02 FRA01 v4;
        protocol udp;
        destination-port 6635;
    then {
        decapsulate mpls-in-udp;
```

Control Plane

Requirements (1)

- → Calculate routing view per Region
 - All machines in a region should have identical routing tables

Requirements (2)

→ Reliable

- Must survive machine failure
- Must support In Service Software Update (ISSU, no it's not a trap)

Requirements (3)

- → Scalable
 - Must support 100+ clients per Region
 - Growing Internet Routing Tables
 - Growing number of Peerings

Requirements (4)

- → Programmable
 - Allow administrative changes to default routing decisions

Getting Routes from Routers

Make BMP Data usable

Getting Routes from Routers (1)

- → BGP Monitoring Protocol (BMP, RFC 7854)
 - Sends all received routes to a monitoring station
 - Notifies monitoring station about peer up/down events
 - Either pre-policy or post-policy
 - We use post-policy

Getting Routes from Routers (2)

- → BIO-Routing Route Information Service (RIS)
 - github.com/bio-routing/bio-rd/cmd/ris
 - Receives BMP messages
 - Tracks per Router/VRF/Peer Adj-RIB-In State
 - Exposes state via gRPC



Getting Routes from Routers (3)



Getting Routes from RIS into SDN Controller

- → Route Information Service (RIS) allows streaming routing information per Router/VRF
- → Uses gRPC Streaming RPC
 - Call ObserveRIB()
 - Reads an (endless) stream of updates
 - RIS sends a state dump initially + updates as they come in via BMP



SDN Controller

Decision Making

Route Controller / SDN Controller (1)

- \rightarrow Written in Go
- → Discovers MPLS Label to Next Hop mapping from IPAM
- → Calculates shortest paths based on BGP data
 - Per Region
 - Per Prefix
 - ♦ BGP Attributes:
 - Local Pref
 - Autonomous System Path
 - MED
 - Origin
 - Internal cost to Next-Hop



Route Controller / SDN Controller (2)

- → Takes Traffic Engineering Input
 - Allows overriding BGP path information
 - To be done automatically
 - Manual action for now

Route Controller / SDN Controller (3)

- → Traffic Engineering Controller is under development
 - Multi-Instance
 - Single leader
 - Takes input from
 - OpenConfig Streaming Telemetry
 - Netflow Collector (tflow2)
 - RIS

Route Controller / SDN Controller (4)

- → Streams Routing Tables to Machines
- → gRPC Streaming RPC
- → New clients receive a full dump
- → Incremental updates sent as route decisions change

Route Attributes:

- Prefix
- Exit Routers Tunnel IP-Address
- MPLS Label
- Weight

Route Agent

Getting Routes into Machines

Route Agent (1)

- \rightarrow Written in Go
- → Makes sure necessary Kernel Modules are loaded
 - ♦ fou
 - mpls_iptunnel
 - mpls_gso

Route Agent (2)

- → Configures Tunnels to Routers
 - Routers are being discovered from Datacenter Inventory Service
- → Maintains a Machines Routing Table
 - Receives Updates from Route Controller
 - Uses Netlink to Replace/Delete Routes in the Linux Kernel

Architecture Overview (Control Plane)



Issues encountered

Go/Netlink issue

- → github.com/vishvananda/netlink
- → Unable to write Multipath Routes with MPLS Encap into the Kernel
- → Encap attribute attached to the wrong object
- → Pull Request waiting for merge

Vendor BMP Issue

- → Router sends incomplete BGP OPEN messages in BMP Peer Up Notifications
- → Only when the peer Router sends exactly 4 Byte-ASN and AddPath capabilities
- → Only when using "allow-from" instead of "neighbor" statement
- → BGP OPEN optional parameters missing

Showing static LSPs briefly as XML output results in invalid XML

- → Only with 100+ LSPs configured
- → JSON output causes segfault





Linux Issues

- → TCP over MPLS Encap Route unusably slow (~70kbyte/s)
 - On a route that made 1,5 Gbps with a non MPLS Route
- → Interface TX drops
- → Random chunks of segments missing
- → Long story short: modprobe mpls_gso

State of Rollout

- → Currently running on internal testing Machines
- → Pending deployment of dedicated SDN Controller Machines
- → Traffic Engineering Controller pending

Thank You!

Questions?